1. Solution. (i)  $\overline{A \cap B} \subseteq \overline{A} \cup \overline{B}$ : Suppose  $x \in \overline{A \cap B}$ . This implies  $x \notin (A \cap B)$ , that is,  $x \notin A$  or  $x \notin B$ . The latter is the same as  $x \in \overline{A}$  or  $x \in \overline{B}$ , that is  $x \in \overline{A} \cup \overline{B}$ .

(ii)  $\overline{A} \cup \overline{B} \subseteq \overline{A \cap B}$ : Suppose  $x \in \overline{A} \cup \overline{B}$ . This implies  $x \notin A$  or  $x \notin B$ . Since x cannot be an element of both A and B, this means  $x \notin A \cap B$ , that is,  $x \in \overline{A \cap B}$ .

## 2. Solution.

- (a) First note that both  $R \odot (S_1 \cap S_2)$  and  $(R \odot S_1) \cap (R \odot S_2)$  are relations from set A to set C. The inclusion is proved by the following chain of implications:  $(a, c) \in R \odot (S_1 \cap S_2)$  implies  $(\exists b \in B) : (a, b) \in R$  and  $(b, c) \in S_1 \cap S_2$  implies  $(\exists b \in B) : (a, b) \in R$  and  $(b, c) \in S_1$  and  $(b, c) \in S_2$  implies  $(\exists b \in B)([(a, b) \in R \text{ and } (b, c) \in S_1]$  and  $[(a, b) \in R \text{ and } (b, c) \in S_2])$  implies  $(\exists b \in B)[(a, b) \in R \text{ and } (b, c) \in S_1]$  and  $(\exists b \in B)[(a, b) \in R \text{ and } (b, c) \in S_2]$  implies  $(\exists b \in B)[(a, b) \in R \text{ and } (b, c) \in S_1]$  and  $(\exists b \in B)[(a, b) \in R \text{ and } (b, c) \in S_2]$  implies  $(a, c) \in R \odot S_1$  and  $(a, c) \in R \odot S_2$  implies  $(a, c) \in (R \odot S_1) \cap (R \odot S_2).$
- (b) Below is one possible example there exist infinitely many others. Let  $A = \{a\}, B = \{b_1, b_2\}, C = \{c\}$  and choose

 $R = \{(a, b_1), (a, b_2)\}, S_1 = \{(b_1, c)\}, S_2 = \{(b_2, c)\}.$ 

Now  $R \odot (S_1 \cap S_2) = \emptyset$  because  $S_1 \cap S_2 = \emptyset$ . On the other hand,  $(a, c) \in (R \odot S_1) \cap (R \odot S_2)$ .

3. Solution. (a) TRUE, (b) FALSE, (c) TRUE(d) TRUE, (e) TRUE, (f) FALSE

## 4. Solution.

- (a) The total number of strings of length 4 is  $10^4$ . There are 10 strings that have 4 identical digits. The number of strings of length 4 that do not have 4 identical digits is 10000 10 = 9990.
- (b) When choosing the first digit we have 10 choices. For a fixed first digit, we have 9 choices for the second digit (to avoid repetition). Similarly, for the 3rd digit we have 8 choices and for the 4th digit 7 choices. The total number of strings that do no repeat a digit is 10 · 9 · 8 · 7 = 5040.

- (c) There are 5 even digits that can be chosen as the last digit. Each of the first 3 digits has 10 choices and using the product rule the total number of choices is  $10 \cdot 10 \cdot 10 \cdot 5 = 5000$ .
- (d) There are 4 ways to choose the positions of the digits 8 (1, 2, 3 or 1, 2, 4, or 1, 3, 4, or 2, 3, 4). The remaining digit can be chosen in 9 ways. Using the product rule the total number of choices if 4 · 9 = 36.
- 5. Solution. We partition elements of A into |B| subsets as follows: two elements  $a_1$  and  $a_2$  belong to the same subset iff  $f(a_1) = f(a_2)$ .

The pigeonhole principle (Theorem 1.2 in the second set of notes,) guarantees that some subset has  $\lceil \frac{|A|}{|B|} \rceil$  elements.

6. Solution. (i) Suppose that f has a right inverse  $g: B \to A$ . For any  $b \in B$  we have

$$b = 1_B(b) = (f \circ g)(b).$$

Since  $(f \circ g)(b) = f(g(b))$ , this means that b is the f image of g(b) and f is onto.

(ii) Conversely suppose that f is onto. For each  $b \in B$ , choose  $a_b \in A$  such that  $f(a_b) = b$ . Define a function  $g: B \to A$  by setting  $g(b) = a_b$  for all  $b \in B$ .

We verify that g is the right inverse of f. For any  $b \in B$  we have

$$(f \circ g)(b) = f(a_b) = b = 1_B(b),$$

that is,  $f \circ g = 1_B$ .