

Computability and Complexity, CISC 462 - Assignment 3 (Fall 2018, K. Salomaa)
Due in lecture 9:30 AM, Monday November 5

1. (a) Find a match for the following instance of the Post Correspondence Problem (PCP).

$$\left\{ \begin{bmatrix} 00 \\ 001 \end{bmatrix}, \begin{bmatrix} 11 \\ 10 \end{bmatrix}, \begin{bmatrix} 011 \\ 1 \end{bmatrix} \right\}$$

- (b) Do the below two PCP instances have a match? Justify your answers.

i. $\left\{ \begin{bmatrix} 1 \\ 101 \end{bmatrix}, \begin{bmatrix} 01 \\ 10 \end{bmatrix}, \begin{bmatrix} 011 \\ 1 \end{bmatrix} \right\}$
 ii. $\left\{ \begin{bmatrix} 01111 \\ 011 \end{bmatrix}, \begin{bmatrix} 101 \\ 11 \end{bmatrix}, \begin{bmatrix} 100 \\ 1000 \end{bmatrix}, \begin{bmatrix} 1010 \\ 110 \end{bmatrix} \right\}$

2. Consider a PCP instance

$$(u_1, v_1), \dots, (u_m, v_m),$$

$u_i, v_i \in \Sigma^*$, $i = 1, \dots, m$. We say that the above instance has a *non-constrained solution* if there exist sequences of integers i_1, \dots, i_s and j_1, \dots, j_r , $r, s \geq 1$ ($i_x, j_y \in \{1, \dots, m\}$), such that

$$u_{i_1} \cdots u_{i_s} = v_{j_1} \cdots v_{j_r}$$

Define

$$L_{ncPCP} = \{ \langle P \rangle \mid P \text{ is an instance of PCP that has a non-constrained solution} \}$$

Show that L_{ncPCP} is decidable.

3. Recall

$$E_{CFG} = \{ \langle G \rangle \mid G \text{ is a context-free grammar and } L(G) = \emptyset \}.$$

- Answer the following questions and **justify** your answers. (“ \leq_m ” denotes mapping reducibility.)

In each case where you answer “yes”, please give a concrete example of the language and the corresponding mapping reduction.

- (a) Does there exist a decidable language A such that $A \leq_m E_{CFG}$?
 (b) Does there exist a decidable language B such that $E_{CFG} \leq_m B$?
 (c) Does there exist an undecidable language C such that $C \leq_m E_{CFG}$?
 (d) Does there exist an undecidable language D such that $E_{CFG} \leq_m D$?

4. Answer each part TRUE or FALSE.

- (a) $n = O(n^2)$
- (b) $n = o(2n)$
- (c) $n^2 = o(n^3)$
- (d) $3^n = 2^{O(n)}$
- (e) $2^{2n} = O(2^n)$
- (f) $2^n = o(2^{n+5})$
- (g) $n^2 = O(n \cdot (\log n)^3)$
- (h) $(\log n)^3 = o(\sqrt{n})$
- (i) $n^2(\log n)^3 = o(n^3)$
- (j) $\frac{1}{n} = o(1)$

5. Define

$$ALL_{\text{DFA}} = \{ \langle A \rangle \mid A \text{ is a DFA and } L(A) = \Sigma^* \}.$$

Show that $ALL_{\text{DFA}} \in \text{P}$.

6. (a) A *triangle* in an undirected graph is a 3-clique and define

$$\text{TRIANGLE} = \{ \langle G \rangle \mid G \text{ is an undirected graph that contains a triangle} \}.$$

Show that $\text{TRIANGLE} \in \text{P}$.

(b) Explain why your algorithm solving TRIANGLE in polynomial time cannot be extended to give a polynomial time algorithm for CLIQUE.

Recall that the language CLIQUE is defined in section 7.3 of the text:

$$\text{CLIQUE} = \{ \langle G, k \rangle \mid G \text{ is an undirected graph with a } k\text{-clique} \}.$$

Regulations on Assignments

- As described on the course homepage, all assignments must be based on *individual work*.
- The assignments are graded according to the correctness, preciseness and elegance of the solutions.
- If, as part of your solution, you rely on results from the textbook you should clearly state which result(s) you are using.
- Each question is worth 10 marks and the assignment is marked out of 60 marks.